

Records/Structures

- Records/structures aggregate heterogeneous data types to be stored and manipulated together
- Records typically declared as distinct types (e.g. type *element* in Scott), particularly for name equivalence
- ML does structural equivalence for records quite well, so no explicit type declaration is necessary:

```
val copper = { name = "Cu", atomic_number = 29, atomic_weight = 63.546 };  
val gold = { name = "Au", atomic_weight = 196.96655, atomic_number = 79 };
```

(* In this case you will have name equivalence *)

```
datatype element = element of { atomic_number : int, atomic_weight : real,  
                               name : string };
```

Record Fields

- Orthogonality: most languages allow records within records; ForTran 90 and Common Lisp do not
- Accessing fields
 - ◇ Dot notation (C, Java, Pascal): `record.nested.field`
 - ◇ Other delimiters (ForTran 90: %): `record%field`
 - * And remember no nested records in ForTran 90
 - ◇ Inverted notation: field first
 - Cobol, Algol 68: `field of nested of record`
 - ML: `#field (#nested record)`
 - Common Lisp: `(field record)`
 - * Again, no nested records here

Records and Memory Layout

- Typically contiguous, in the declared field order
 - ◇ Alignment issues: the classic time vs. space tradeoff
 - ◇ One way to optimize is to rearrange the fields internally, but this will be an issue in systems programming where record layouts are supposed to mirror low-level bit fields
- Pascal: explicit *packed* keyword to designate preference to save space rather than time
- Memory layout also affects record comparison and assignment
 - ◇ *Assignment*: bit-for-bit transfer (block copy)
 - ◇ *Comparison*: bit-for-bit compare (block compare), maybe, but...how about the holes due to field alignment?
 - Can be tricky — so Pascal and C ditch the whole idea

Scoping Records

- Multiple activities on the fields of the same record can result in lengthy code, particularly for “deep” records
- Pascal: *with* keyword sets up a local scope; identifiers for fields can then be addressed “by themselves”
 - ◇ How to manipulate two records of the same type within a *with* block?
 - ◇ How about standalone variables with the same name as record fields?
- C/C++: instead of *with*, setup arbitrary pointers whose base type is the record, then use \rightarrow notation
- Object-oriented languages: functions/methods within classes set up this scope for member/instance variables of those classes

Variants/Unions

- Two or more alternative fields or collections of fields, only one of which is valid at any given time

- ◆ In other words, “memory recycling”

- Tag/discriminant: record field that determines how to interpret the shared memory

- ◆ Tag can be named or anonymous

- For Tran, Algol: *equivalence* — memory sharing for multiple variables, not just within a record
- “Integration” — does the variant section have a name of its own (e.g., in C, it does; in Pascal, it does not)

- Semantic checks galore!

- ◆ When to use which alternative representation

- Check the tag/discriminant?
- What if the tag is anonymous?
 - Algol 68: assignments change “which” alternative is current
- Compile time or runtime? Sometimes it is impossible to do compile time — implies additional runtime code

- ◆ What about initialization — say we change the tag value on the fly?

- Due to these issues, latest descendants of certain languages have dropped variants/unions: Modula-3, Java
- Primary use today: systems programming, where bit fields may be interpreted in more than one way

- ◆ Color may be a 4-byte red/green/blue/alpha structure or a single unique 32-bit integer

Arrays

- Homogeneous mappings from a domain of *index* values to a range of *component* or *element* values
- Basic version restricts index to integers (traceable to origins of how arrays are implemented) — C and descendants force zero-based indexing
- Some languages allow subranges or enumeration types
- Most generalized version can map any type to any other type — these *associative arrays* are sufficiently different in terms of implementation that they are named differently, usually *maps* (Java, C++)

Array Elements

- Array elements originally scalar (ForTran, BASIC)
- Now, elements can be of any type, even other arrays
- Access by subscript, delimited by parentheses (ForTran, Ada) or square brackets (Pascal, C/C++, Java)
 - ◊ Square bracket advantage: avoids parenthesis overloading as function arg delimiters
 - ◊ ...unless you explicitly want to think of array accesses as a form of function call! (Ada)
- Arrays are typically declared with a fixed size
- Arrays of arrays are *multidimensional* — i.e., a single basic array counts as one dimension (geometric)
- Pointers and arrays in C/C++ — a special case

Array Slices and Operations

- Array slices are subsets of a larger array; heavily supported in Fortran 90, supported for 1-dimensional arrays in Ada only, largely unsupported elsewhere
- Array operations: most of the time, purely access and assignment only (essentially, get and put)
 - ◆ Ada and Fortran 90 allow comparisons and operations over every element of an array (Ada: boolean operators; Fortran 90: arithmetic, mathematical functions)
 - ◆ Other languages require explicit looping over elements of an array
 - ◆ Array iterators in newer languages simplify things somewhat, though they may never reach the succinctness of Fortran 90 array addition:

$$C = A + B$$

Dimensions, Bounds, Allocation

- Array implementations can be distinguished by their “life cycles,” and life cycle implies binding
- For arrays: when to bind to memory, and when to bind their shape (size, dimensions, size per dimension)
 - ◆ *global lifetime, static shape*: available through program life, fixed shape at compile time
 - ◆ *local lifetime, static shape*: allocated at runtime (local variable), fixed shape at compile time
 - ◆ *local lifetime, shape bound at elaboration time*: allocated at runtime, at a fixed shape once allocated, but that shape is not known until runtime
 - ◆ *arbitrary lifetime, shape bound at elaboration time*: essentially a heap-allocated array; fixed shape, but known only at runtime
 - ◆ *arbitrary lifetime, dynamic shape*: heap-allocated array whose shape can change at any time

Arrays and Memory

- The “first array” — contiguous locations in memory
- For arrays of composite types, memory alignment is an issue just as with records
 - ◇ Pascal *packed* keyword applies to both arrays and records, to similar effect
- How about multidimensional arrays?
 - ◇ *Row-major vs. column-major*: is the next element in memory the next array element along the same dimension, or the equivalent array element along the next dimension?
 - ◇ *Key difference*: easier to view a multidimensional array as nested arrays of arrays, since the nesting remains contiguous in memory
 - ◇ *Performance issues*: arrangement may depend on how array elements are iterated
 - ◇ *Alternative representation*: instead of contiguous blocks of the element type, an array can be contiguous pointers to blocks of that type even for single dimensions

- The general approach for contiguous layout, given a known start location in memory, is an accumulated sum of products of indices along each dimension and the size of the element type; pointer-based layout is similar, using the size of a pointer instead of the size of the element type
- But — can we perform this at compile-time or runtime? In other words, can we pre-calculate the memory offset of an array element at compile time?
 - ◇ In other words, can $A[i, j, k]$ be automatically compiled as some $memory[offset]$?
 - ◇ Depends on knowledge of array shape
 - ◇ Also depends on how array is allocated (statically, on the stack, on the heap) — this determines knowledge of the start location

Bounds Checking

- Accesses outside of array bounds are always semantic errors — static vs. dynamic depends on the language
- C/C++ — because arrays and pointers are essentially variations on the same theme, must wait for “Segmentation fault” or similar-sounding errors to realize that you are out-of-bounds...and sometimes you never find out
- In other languages (such as Java), an array “knows” how long it is; more generally, an array “knows” its shape
 - ◇ Shape is stored in a dope vector for the array
 - ◇ Still, does not completely eliminate dynamic semantic checks

Strings

Structurally, strings are “just” one-dimensional arrays of characters

- In many languages, that’s all they are (Pascal, C)
- Other languages have a specific *string* type (Java, C++, ML, JavaScript)
- Either way, even in languages where strings are “just character arrays,” they frequently get special handling (Pascal: string literals, C: complete string library)

Strings as the Ultimate “Special Case”

Strings are a special case, and a *sufficiently* special case:

- Broad applicability — virtually every program needs to manipulate strings
- Strings have a stable, consistent structure (one-dimensional array of bytes, pre-Unicode; 16 bits per character for Unicode)
- Makes implementation of string manipulation easier than implementing the generalized array manipulation case, allowing optimizations

String Operations

- String literals: specifying string values directly in code
 - ◇ Support for “escape characters” when content is not typable: `\n`, `\t`, `\u` (Unicode escape)
- String functions (length, concatenation, comparison) — note how they are conceptually generalizable to arrays of any type, but are difficult to implement that way
- C: string functions are thin wrappers on generalized character pointer manipulation
- Java: String is a distinct, full-fledged class
- JavaScript: strings are a specific type, and also serve as a “bridge type” for other types

Sets

- Unordered collection of values for a particular element type, with values appearing only once in the collection
 - ◇ Key operations are the same as sets in math: union, intersection, difference
- In the huge, general case, sets may be viewed as specialized arrays (or maps)
 - ◇ Remove access by order
 - ◇ Check for duplicate elements
- In special cases, sets can be implemented as a bit field and thus can be very fast
 - ◇ One bit per element in the type's domain
 - ◇ Zero if that element is in the set, one if not
 - ◇ A zero bit field is the empty set
- This approach genuinely reflects “no duplicates” semantics of sets, but is impractical when the set's base type is very large
- Thus, languages with this type of set frequently limit the base types to small ones, typically 128 or 256 elements
- Hash tables can serve as an alternative implementation when base types are large

Pointers/Recursive Types

- Recursive types are types whose definitions include references to themselves (or to types that eventually reference back in a cycle)
- Useful in many data structures: linked lists, trees, graphs
- Type recursion leads to some form of reference
 - ◇ Implicit reference: Lisp, ML, Clu, Java
 - ◇ Explicit reference (therefore pointer type): C, Pascal, Ada
- Pointers are *not* addresses
 - ◇ Addresses are direct locations in memory
 - ◇ Pointers are references to a separate object without being the object itself — they may be addresses, but not necessarily

Recursive Type Implementation

- Reference model
 - ◇ All symbols are references, but can be accessed without additional syntax
 - ◇ Thus, recursive types just restate the type
 - ◇ An assignment copies a reference, not its value
- Value model
 - ◇ Special pointer type explicitly differentiates a reference
 - ◇ Recursive types are defined through pointers to that type
 - ◇ Subfields/elements of the recursive type must be dereferenced through the pointer using specific syntax (Pascal: ^; C: *, ->)
 - ◇ Assignments always copy values — note how this is still consistent with pointer types: assigning one pointer to another does copy that pointer's value into the other variable

Pointers and Allocation

In relation to how objects are allocated in a program, pointers can be:

- A mechanism for heap allocation of objects
- A mechanism for referring to any other value in a program (individual variables, record fields, array elements, subroutines)
 - ◇ Implies availability of an “reference to” operator (C/C++: `&`; ML: *ref*)
 - ◇ Pointers are thus a form of binding too
- Depending on the language, pointers may be exclusively for heap use; sometimes both
 - ◇ Pascal, Ada 83, Modula-3: heap only
 - ◇ PL/I, Algol 68, C, C++, Ada 95: heap or reference
- When pointers refer to the heap, the natural follow-up question is how to reclaim heap space
 - ◇ Manual reclamation: C, C++, Pascal, Modula-2
 - ◇ Automatic reclamation (*garbage collection*): Lisp, ML, Modula-3, Ada, Smalltalk, Java
 - ◇ Resource/performance constraints once preferred manual reclamation; these days, the additional cost of garbage collection has been offset by the cost of finding and fixing memory allocation bugs

Case Study: Pointers and Arrays in C

- C views arrays and pointers as interchangeable
- An array is a pointer; multidimensional arrays are n -level pointers to pointers
- Dereferencing an array is the same as “moving” a pointer down the array an *index* number of times
- Pointer arithmetic allows numerical manipulation of pointers; must take into account the size of the pointer’s base type
- Array layout becomes an issue

Pointer Management

- *Dangling references* — pointers whose *referents* are no longer bound (heap deallocation, call stack)
- Converse of dangling references are *memory leaks* — referents that are bound in memory but have nothing pointing to them
- Multiple techniques for helping with dangling references: *tombstones*, *lock and key*, *destructor* methods or functions (explicit clean-up code)
- Garbage collection techniques: *reference counts*, *mark-and-sweep* (heap traversal)

Lists

- Another collection type of a specific base element type, internally implemented as a “lower level” collection such as a specific array or a linked list
- Presented for recursive access as a *head* and a *tail* (Lisp, ML); otherwise, behaves like array-like collections
 - ◇ Head has base element type
 - ◇ Tail is another list
- Operations for list concatenation
- Special literal for the empty list
- Iteration-controlled loops are applicable

Streams

- Arrays with “state”
 - ◇ “Current element,” whether or not we are at the end of a stream
 - ◇ Explicit traversal “up” or “down” the array
- Two directions
 - ◇ Input stream — stream from which we get values
 - ◇ Output stream — stream to which we put values
- Originated with the specific needs of file I/O, but since generalized (Java, C++)
 - ◇ Files remain, now modeled as streams of various types (bytes, characters, strings, objects)
 - ◇ Blocks of memory (byte arrays, strings) can be viewed as streams
 - ◇ Database constructs: tables, query results